

# 6.172 Performance Engineering of Software Systems

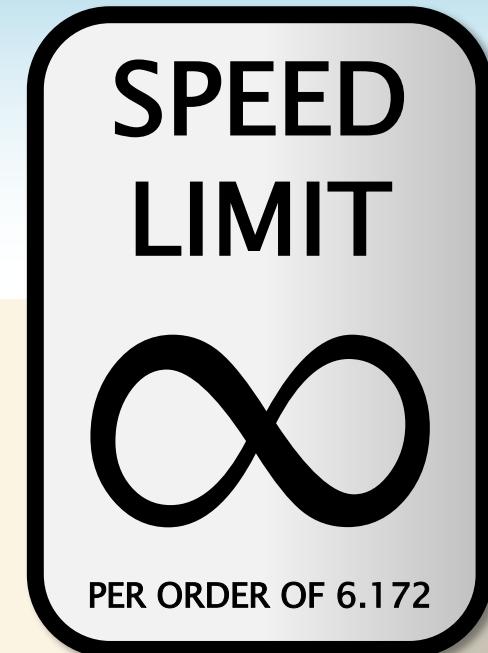


REVIEW

## Divide-and-Conquer Recurrences

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# The Master Method

The **master method** for solving divide-and-conquer recurrences applies to recurrences of the form\*

$$T(n) = aT(n/b) + f(n),$$

where  $a \geq 1$ ,  $b > 1$ , and  $f$  is asymptotically positive.

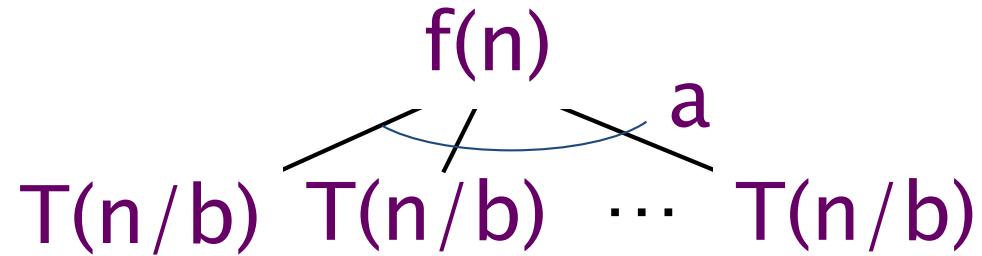
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\*The unstated base case is  $T(n) = \Theta(1)$  for all  $n \leq n_0$  for a sufficiently large constant  $n_0$ .

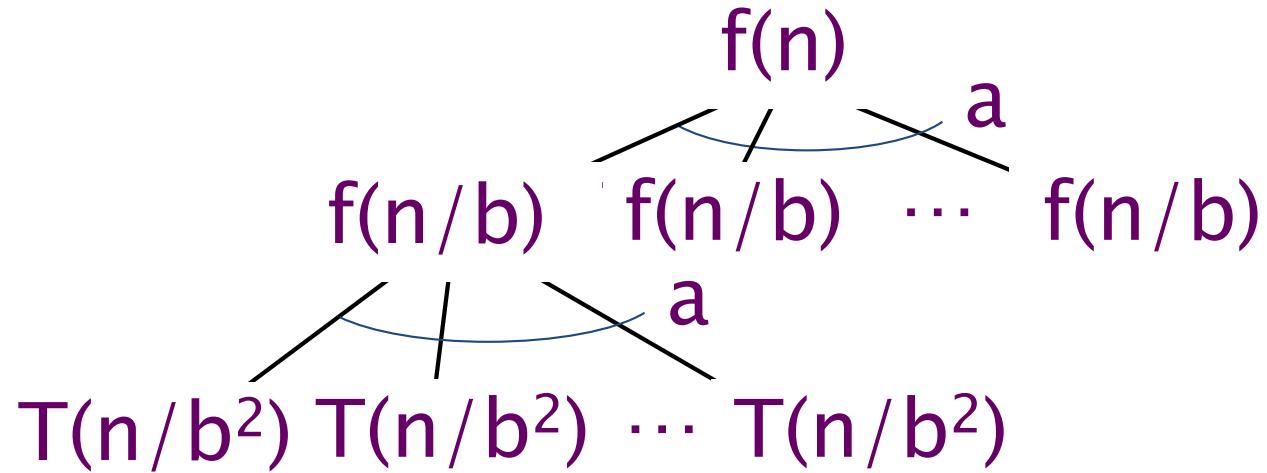
# Recursion Tree: $T(n) = aT(n/b) + f(n)$

$T(n)$

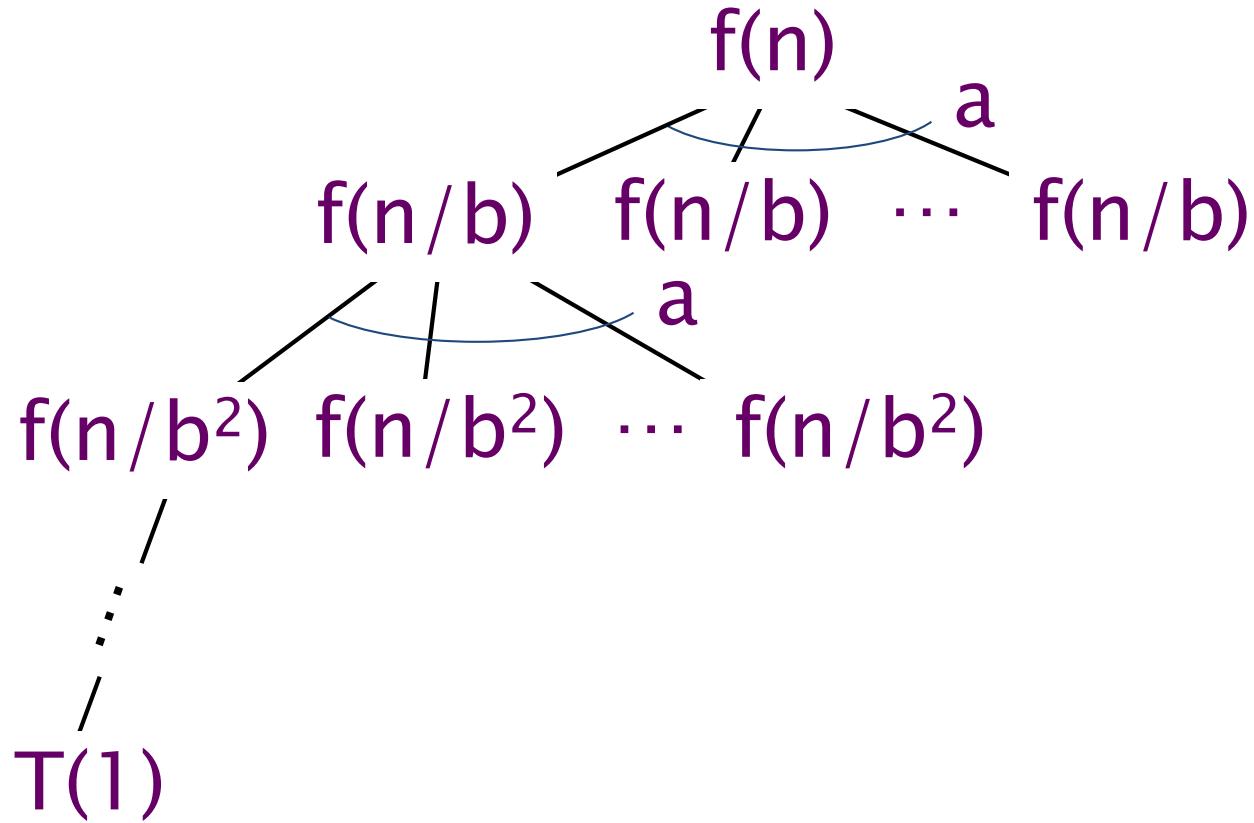
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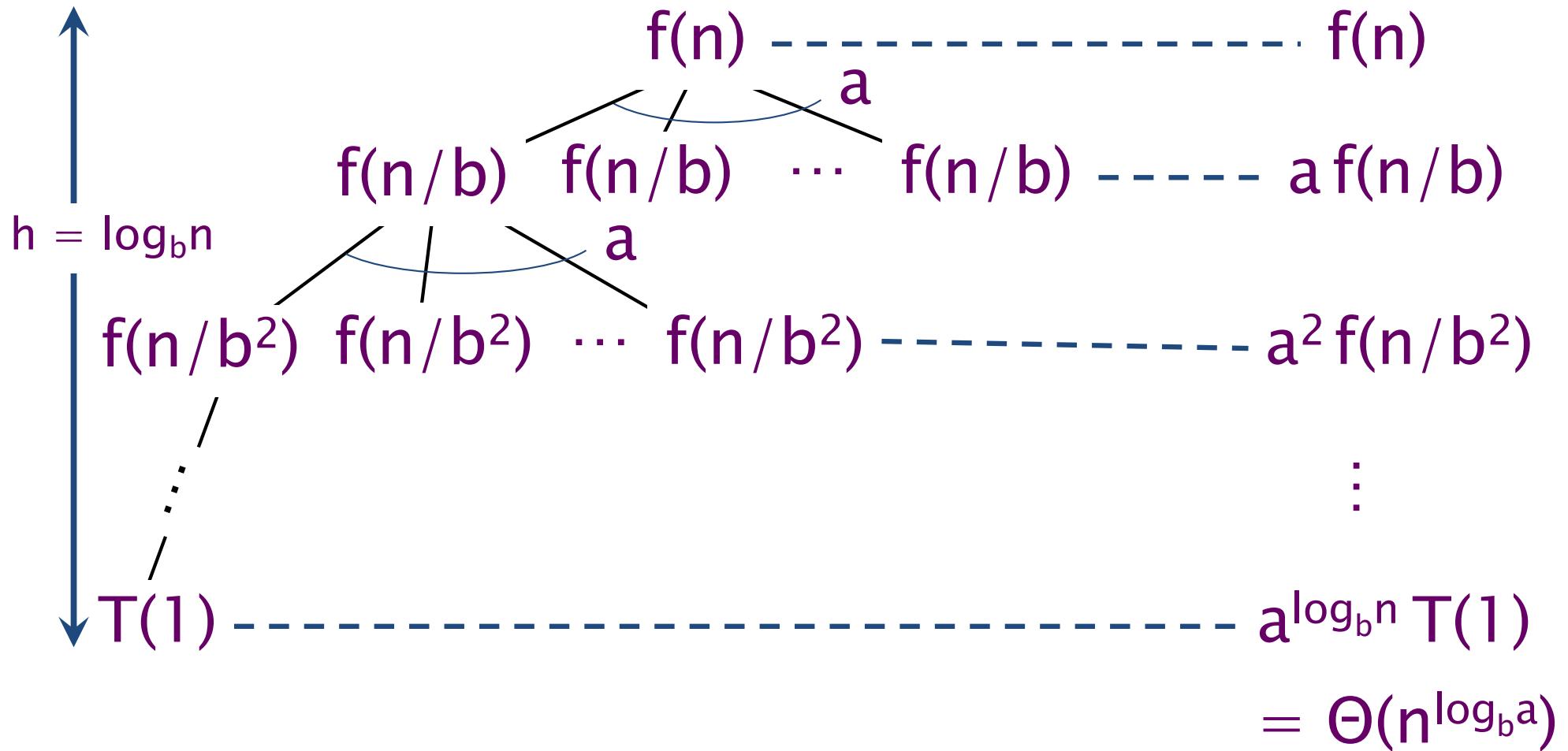
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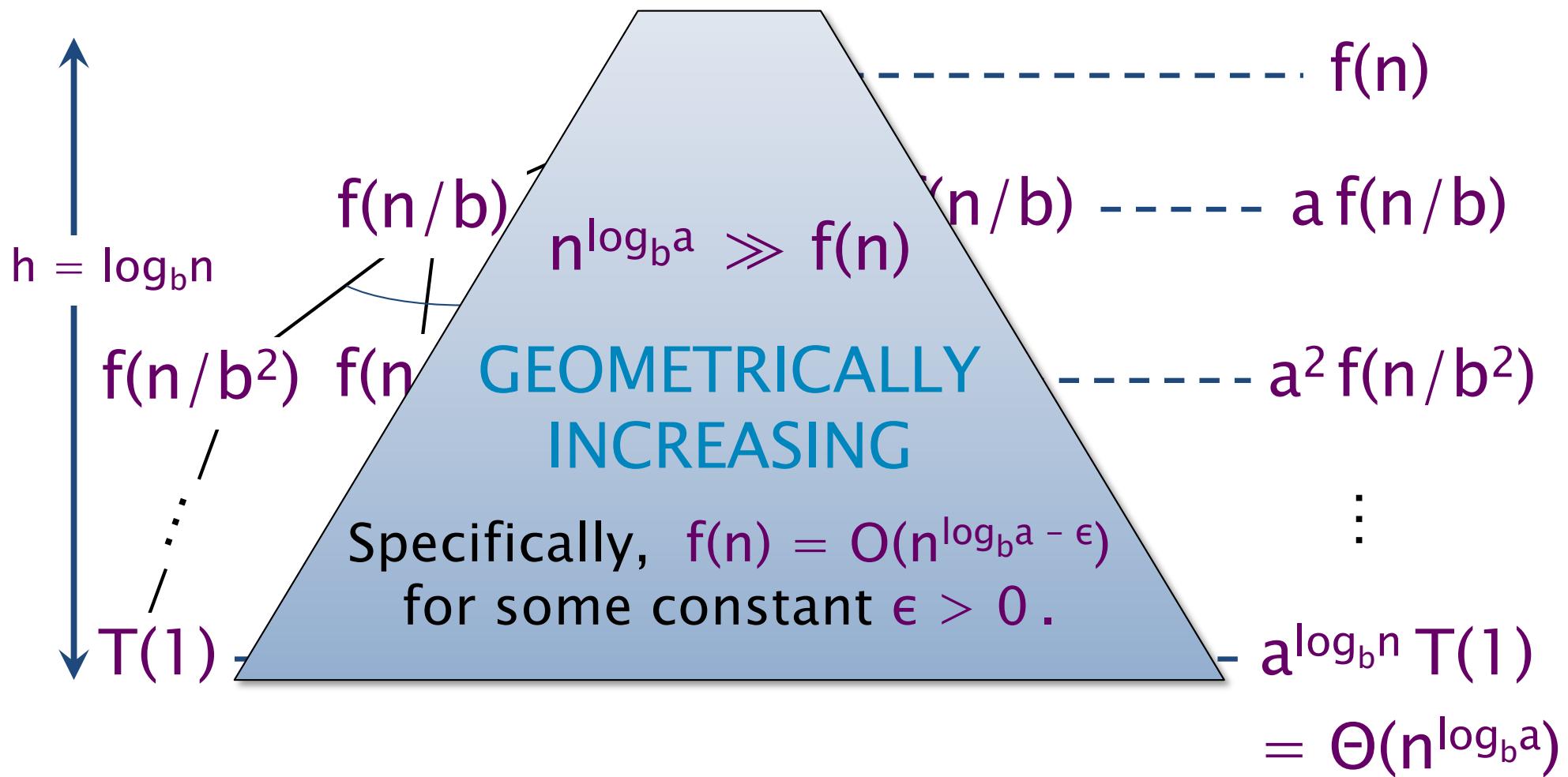


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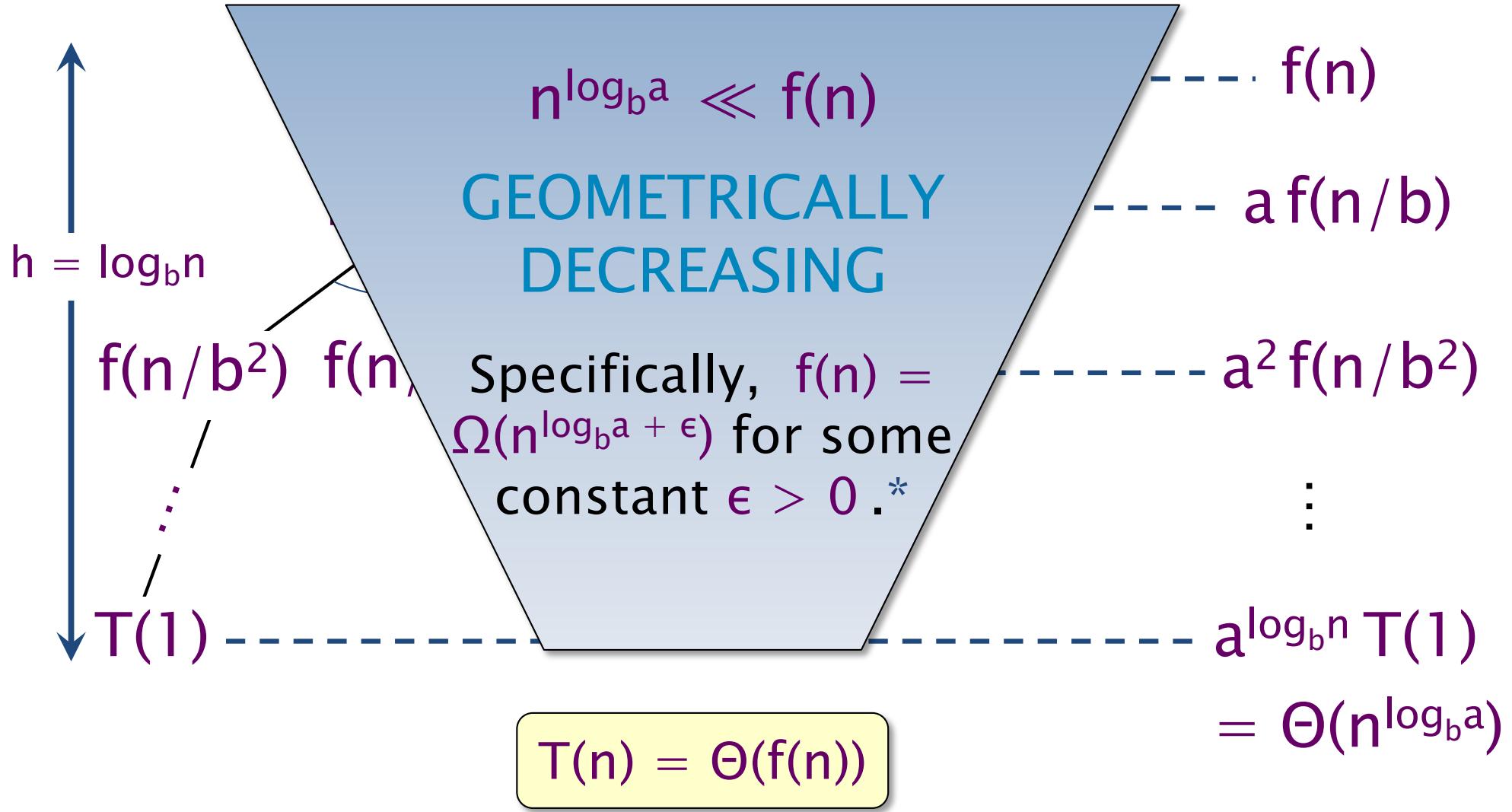


IDEA: Compare  $n^{\log_b a}$  with  $f(n)$ .

# Master Method — CASE I

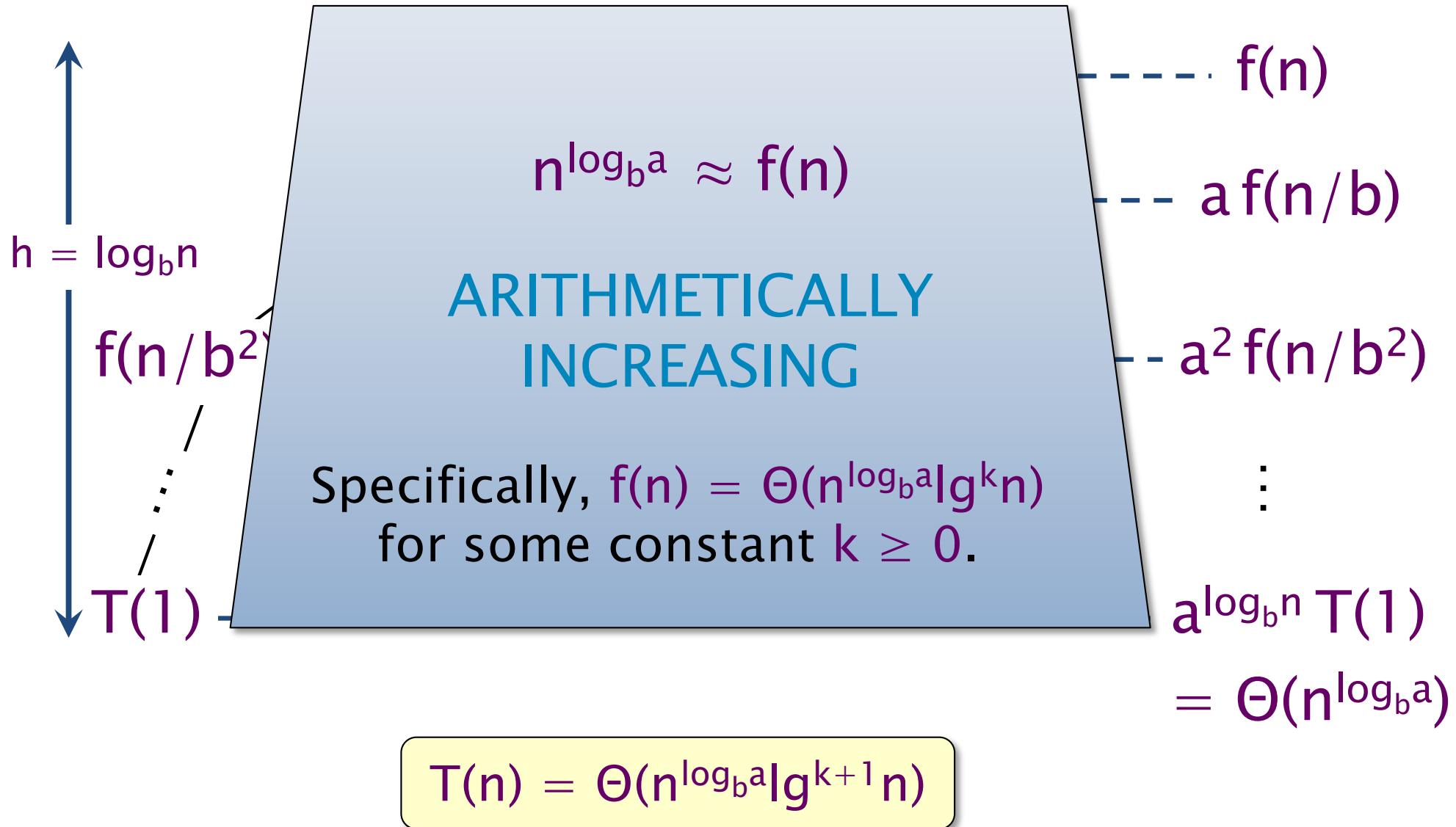


# Master Method — CASE 3



\*and  $f(n)$  satisfies the **regularity condition** that  $a f(n/b) \leq c f(n)$  for some constant  $c < 1$ .

# Master Method — CASE 2



# Master-Method Cheat Sheet

Solve

$$T(n) = aT(n/b) + f(n) ,$$

where  $a \geq 1$  and  $b > 1$ .

**CASE 1:**  $f(n) = O(n^{\log_b a - \epsilon})$ , constant  $\epsilon > 0$   
 $\Rightarrow T(n) = \Theta(n^{\log_b a})$  .

**CASE 2:**  $f(n) = \Theta(n^{\log_b a} \lg^k n)$ , constant  $k \geq 0$   
 $\Rightarrow T(n) = \Theta(n^{\log_b a} \lg^{k+1} n)$  .

**CASE 3:**  $f(n) = \Omega(n^{\log_b a + \epsilon})$ , constant  $\epsilon > 0$   
(and regularity condition)  
 $\Rightarrow T(n) = \Theta(f(n))$  .

<https://tinyurl.com/mm-cheat>

# Master Method Quiz

- $T(n) = 4T(n/2) + n$   
 $n^{\log_b a} = n^2 \gg n \Rightarrow \text{CASE 1: } T(n) = \Theta(n^2).$
- $T(n) = 4T(n/2) + n^2$   
 $n^{\log_b a} = n^2 = n^2 \lg^0 n \Rightarrow \text{CASE 2: } T(n) = \Theta(n^2 \lg n).$
- $T(n) = 4T(n/2) + n^3$   
 $n^{\log_b a} = n^2 \ll n^3 \Rightarrow \text{CASE 3: } T(n) = \Theta(n^3).$
- $T(n) = 4T(n/2) + n^2/\lg n$   
Master method does not apply!  
Answer is  $T(n) = \Theta(n^2 \lg \lg n)$ . (Prove by substitution.)

More general (but more mathematically sophisticated) solution: *Akra-Bazzi method.*