Software Performance Engineering



Recitation 1.6

Sophia Sun Tuesday, October 7, 2025

Some due dates

- Project 1 Final is due this Thursday!
- You can use late days for your writeup but not your code submission.
- Make sure to run the correctness check first!
- Reach tier 35 to get at least a B grade for final.

Homework 4 is due next Monday (October 13).

The Final Writeup

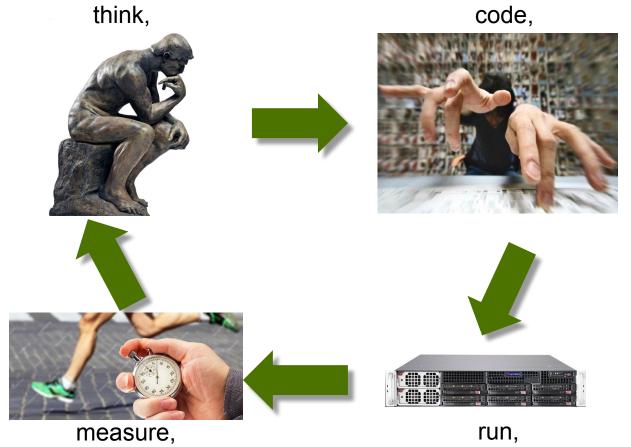
 Your final write-up can be an update of your beta write-up. In addition to the beta, it should include the following:

- An overview of changes you made to your beta release and what motivated you to do them
 - surprised by performance ranking
 - ideas conceived before the beta deadline but ran out of time implementing them
- Any updates to the acknowledgment, including specifically which of classmates' beta codes may have inspired you.



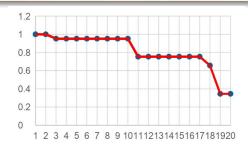
PERFORMANCE ENGINEERING

Performance Engineering



Basic Performance-Engineering Workflow

- 1. Measure the performance of Program A.
- 2. Make a change to Program A to produce a hopefully faster Program A'.
- 3. Measure the performance of Program A'.
- 4. If A' beats A, set A = A'.
- 5. If A is still not fast enough, go to Step 2.





If you can't measure performance **reliably**, it is hard to make many small changes that add up



Tools to Measure Software Performance

How Much to Measure

- Measure the whole program
 - □ e.g., /usr/bin/time
 - □ e.g., perf stat, cachegrind, strace
- Measure just the part of the program we care about
 - Insert timing calls in the program
 - e.g., clock_gettime(), gettimeofday(), rdtsc()
- Create a profile of the program
 - e.g., gdb, gprof
 - use perf_record/report, based on HW counters that OS and HW support
 - Using sampling or instrumentation

/usr/bin/time

The time command can measure elapsed time, user time, and system time for an entire program.

```
$ /usr/bin/time my-program arg1 arg2
real 0m3.502s
user 0m0.023s
sys 0m0.005s
```

What does that mean?

- real is wall-clock time (elapsed time).
- user is the amount of processor time spent in user-mode code (outside the kernel) within the process.
- sys is the amount of processor time spent in the kernel within the process

clock_gettime(CLOCK_MONOTONIC, ...)

- Typically, clock_gettime(CLOCK_MONOTONIC, ...) is fast
 roughly 80ns about 10² faster than an ordinary system call
- clock_gettime(CLOCK_MONOTONIC, ...) has nice guarantees
 it never runs backwards
- But clock_gettime(CLOCK_MONOTONIC, ...) isn't always fast

rdtsc()

x86 processors provide a time-stamp counter (TSC) in hardware.

You can read TSC as follows:

For older compilers

For newer compilers

```
__builtin_readcyclecounter();
```

- The time returned is "clock cycles since boot."
- rdtsc() runs in about 32ns.

Issues with TSC

Use rdtsc() with caution!

- May give different answers on different cores on the same machine
- TSC can appear to run backwards, due to process migration
- Not all cycles take the same amount of time!
 - Modern processors change clock frequencies dynamically, e.g. DVFS, Turbo Boost
 - Processors reduce their clock frequency for AVX, AVX2, and AVX512 instructions
- Converting clock cycles to seconds can be tricky

Program Instrumentation

We can make a profiler by instrumenting the program.

```
static vec_t vec_add(vec_t a, vec_t b) {
                                                     Instrumentation
 clock_gettime(CLOCK_MONOTONIC, &start);
 vec_t sum = { a.x + b.x, a.y + b.y };
 clock_gettime(CLOCK_MONOTONIC, &end);
                                                     Instrumentation
 report_time("vec_add", &start, &end);
 return sum;
static vec_t vec_scale(vec_t v, double a) {
                                                     Instrumentation
 clock_gettime(CLOCK_MONOTONIC, &start);-
 vec_t scaled = { v.x * a, v.y * a };
                                                     Instrumentation
 clock_gettime(CLOCK_MONOTONIC, &end);
 report_time("vec_scale", &start, &end);
 return scaled;
```

Caution: Watch out for probe effect, where program instrumentation alters the program's behavior in unintended ways.

Profiling by Sampling

IDEA: Interrupt the running program at regular intervals, and look at the stack each time to determine which functions are usually being executed.

- pmprof, gprof, perf record, and gperftools automate this strategy to provide profile information for all functions
- This approach is not accurate if it doesn't obtain enough samples. (gprof samples only 100 times per second.)
- You can do this yourself!
 - "Poor Man's Profiler": Run your program under gdb, and type control-C, but at random intervals

Simulation

Simulators can deliver accurate and repeatable perf numbers

- □ **e**.**g**. cachegrind, gem5
- For deterministic programs, you need only run the simulator once
- Simulators are robust to probe effect
 - If want a particular metric, you can go in and collect it without perturbing the simulation
- But they often run much slower than real time
- Simulators don't capture all relevant performance phenomena
 - e.g. cachegrind does not model all memory-system features, e.g., prefetching

Performance Surrogates

What could we measure as a surrogate for time?

- Work: Number of executed instructions
 - Using hardware counters (on systems that support them) or program instrumentation
- Processor cycles
 - Using rdtsc()
- Memory accesses, or cache hits and misses
 - Using hardware counters
 - cachegrind simulation (gives repeatable numbers but slow)
- Span
 - Using Cilkscale



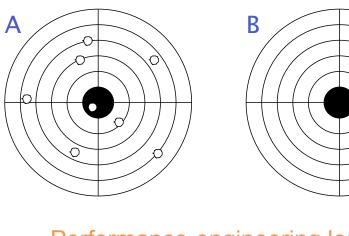
QUIESCING SYSTEMS

Genichi Taguchi and Quality

Question: If you were an Olympic pistol coach, which shooter would you recruit

for your team?

Answer: B, because you just need to teach B to shoot lower and to the left.



Performance-engineering lesson

If you can reduce variability, you can compensate for systematic and random measurement errors.

Sources of Variability

- Daemons and background jobs
- Interrupts
- Code and data alignment
- System calls
- Operating-system process scheduling
- Thread placement

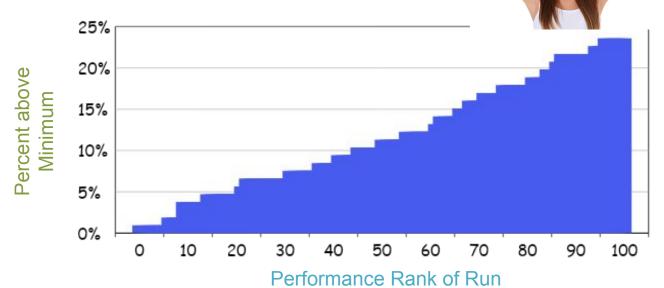
- Runtime scheduler
- DVFS and Turbo Boost
- Network traffic
- Multitenancy
- Virtualization
- Hyperthreading



Unquiesced System

Experiment

- Cilk program to count the primes in an interval
- AWS c4 instance (18 cores)
- · 2-way hyperthreading on, Turbo Boost on
- 18 Cilk workers
- 100 runs, each about 1 second



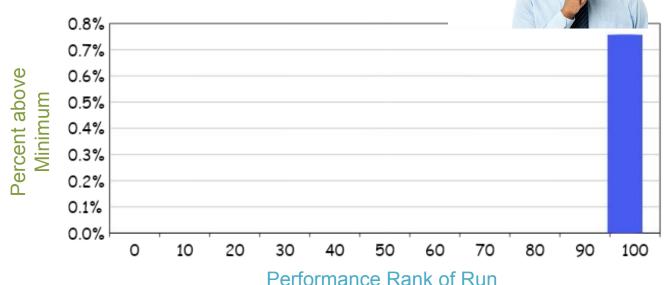
Quiesced System

Experiment

Cilk program to count the primes in an interval

AWS c4 instance (18 cores)

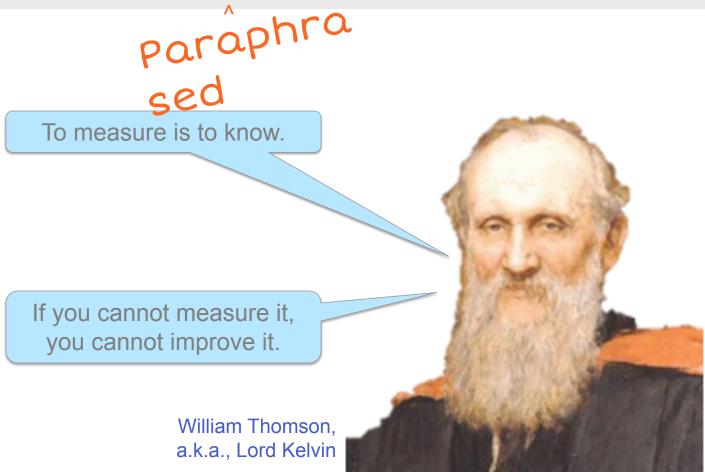
- 2-way hyperthreading off, Turbo Boost off
- 18 Cilk workers
- 100 runs, each about 1 second



Quiescing the System

- Minimize the number of other jobs running
 - Shut down daemons and cron jobs
 - Disconnect the network
 - □ Don't fiddle with the mouse!
 - For serial jobs, don't run on core 0, where many interrupt handlers are usually run, see /proc/interrupts
- Use Linux CPU frequency governor to control DVFS and Turbo Boost
- Use taskset to pin Cilk workers to cores or hardware threads and avoid hyperthreading
- Many of these mitigations have been done for you in telerun

Wise Words from a Giant of Science





CODING TIME BEFORE FINAL DUE!

If you still haven't, think about these...

- How did you set up your temp buffer for loading and storing your blocks? How well does it utilize your cache?
- Any possibility that you can try out some different buffer size? How do you allocate your buffer if you want to load a different block size?
- How can you overlap some direct memory access time with the actual rotation time? (Hint: think about prefetching)
- In what order are you accessing the matrix right now when reading and writing blocks? How can you access your matrix in a cache-friendly way?